An Optimum Method for Waiting Time Distribution: M/D/C

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Abstract: A probabilistic investigation determines the ideal and unequivocal articulation for the holding up time conveyance of the M/D/C lining framework. In this paper, the strategy for numerical difficulties isn't required for high traffic intensities. By fulfilling Erlang's fundamental condition for the M/D/C line, the outcomes can be demonstrated unequivocal. Basic recipe for the holding up time conveyance driving by an option probabilistic approach, which is numerically steady for all $\rho < 1$. With the aiding of holding up time appropriation and Erlang's indispensable condition, we demonstrated a few outcomes for M/D/C dispersion at the base of quality of clients, size of the line.

Key words: Arbitrary customers, waiting time distribution, Markov chain, Erlang's integral equation.

M/D/C system: In the beginning of this century it has been already studied. This system is one of the best classical queuing models. In this model we suppose **c** identical servers, serving each customer during a constant time D. Arrival of customers is according to a Poisson process with rate λ .Now waiting time distribution of the M/D/C queue by Erlang.

$$F(y) = \int_{0}^{\infty} F\left(x + y - D\right) \frac{\lambda^{C} x^{c-1}}{(c-1)!} e^{-\lambda x} dx, \quad y \ge 0 \& c$$

.....(1)
Where, $\frac{\lambda^{C} x^{c-1}}{(c-1)!} e^{-\lambda x} dx = \Pr{obability}$

Two a few clients are in holding up line. The condition (1) is gotten by looking at the holding up time of the clients An and C, who will touch base amongst x and x+dx time. Erlang understood that for c>1.it would scarcely prompt an express scientific arrangement. Crommelin additionally determined a general articulation for the holding up time conveyance of the M/D/C line for all c ϵ N. A recursion plot in view of Crommelin's contention is depicted in Tijms to get around the issue round off mistakes. Be that as it may, for expanding c and ρ this recursion plan will at last be hampered by round off blunders.

Strength of customers: Let there are i customers holding at time t & the probability is $P_i(t)$. So, customers present at time (t + D) either arrived during the time interval (t, t + D].

i.e.
$$P_i(t+D) = \sum_{j=0}^{c} P_j(t) \frac{\lambda D}{i!} e^{-\lambda D} + \sum_{j=c+1}^{i+c} P_j(t) \frac{(\lambda D)^{i+c-j}}{i+c-j!} e^{-\lambda D}, \quad t \in R, i \in N_0$$

.....(2)

If

Where $P_i =$ Stationary distribution.

Queue size: Let A be the landing time of nth client after t = 0, Sn is the beginning time of the administration of the nth client, Wn is the

holding up time of the nth client, qi+(t) is the likelihood of finding a line of length i at time t+, Lq+(t) is the line length (size of line) at time t+, qi is stationary likelihood and , we will demonstrate the accompanying articulation for the stationary probabilities of having I holding up clients in the line instantly after some discretionary administration begin.

$$W_i = \lim_{t \to \infty} P\{L_q(t) = i\} = q_{i,} \quad i \in N_0$$

Proof:

Give S_n a chance to be the administration beginning of nth client while i clients left in the line and $q_i+(S_n)$ is a likelihood with this administration. Amid benefit interim is the likelihood of j fresh introductions. The server will begin serving the (n + e)th client promptly in the wake of completing the administration of the nth client while So size of the line = (i + j - c)

In the event that the client has not yet landed at time so promptly in the wake of completing the administration of the nth client. Consolidating the two cases we infer that Therefore the likelihood of having I holding up clients in the line quickly after the administration beginning of the client.

$$q_{o}^{+}(S_{n+c}) = \sum_{j=0}^{c} q_{j}^{+}(S_{n}) \sum_{m=0}^{c-j} \frac{(\lambda D)^{m}}{m!} e^{-\lambda D}$$
$$q_{i}^{+}(S_{n+c}) = \sum_{j=0}^{i+c} q_{j}^{+}(S_{n}) \frac{(\lambda D)^{i+c-j}}{(i+c-j)!} e^{-\lambda D}$$

If
$$n \to \infty$$
, then the stationary distribution,

$$w_{0} = \sum_{j=0}^{c} w_{j} \sum_{m=0}^{c-j} \frac{(\lambda D)^{m}}{m!} e^{-\lambda D}$$
$$w_{i} = \sum_{j=0}^{i+c} w_{j} \frac{(\lambda D)^{i+c-j}}{(i+c-j)!} e^{-\lambda D}$$
While i > 0
.....(5)

 $q_i \forall i \ge 0$

Hence w_i equals

$$w_0 = \sum_{j=0}^{c} w_j \sum_{m=0}^{c-j} \frac{(\lambda D)^m}{m!} e^{-\lambda D}$$

$$w_i = \sum_{j=0}^{i+c} w_j \frac{(\lambda D)^{i+c-j}}{(i+c-j)!} e^{-\lambda D}$$

The waiting time Distribution: Our first target will be to decide time " W_n " of the nth clients to touch base after t = 0. Client is called checked client and server is called stamped server. Watch that stamped client will be the K_{th} client to be served by the checked server from time moment onwards. Give u a chance to be some positive time slip by < D. On the off chance that the checked client arrives no sooner than the stamped server has been serving the client for at any rate u time units on the entry moment An. Along these lines the checked client will locate the stamped server with a measure of incomplete work < kD-u, suggesting . In the second case, if the checked client lands before.

In the event that at $L_q^+(S_{n-kc})$, the line length quickly after S_{n-kc} . The line contains kc or more clients, the checked client has just entered the line and is holding up there in (kc)th position, suggesting A < S_{n-kc} . With likelihood Now stationary time distribution by letting $n \rightarrow \infty$,

$$P\{W \le kD - u\} = \sum_{i=0}^{kc-1} q_i \sum_{j=0}^{kc-i-1} \frac{(\lambda u)^j}{j!} e^{-\lambda u} = e^{-\lambda u} \sum_{j=0}^{kc-1} \frac{(\lambda u)^j}{j!} \sum_{i=0}^{kc-i-1} q_i$$
......(7)

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Substituting kD-u=x, then the waiting time distribution

$$P\{W \le x\} = e^{-\lambda(kD-x)} \sum_{j=0}^{kc-1} Q_{kc-j-1} \frac{\lambda^j (kD-x)^j}{j!}$$

..... (8)
While $(k-1)D \le x < kD$

Erlang's integral Equation: In order to complete the circle, this section presents an explicit proof which will make use the following lemma:

$$Q_n = \sum_{i=0}^{n+c} Q_{n+c-i} \frac{\left(\lambda D\right)^i}{i!} e^{-\lambda D} \quad \forall \quad n \in N_0$$

Proof: Let no more than n customers are in queue at time and $Q_n(t)$ is the probability of the queue

$$Q_n(t+D) = \sum_{i=0}^{n+c} Q_{n+c-i} \frac{\left(\lambda D\right)^i}{i!} e^{-\lambda D} \quad \forall \quad n \in N_0$$
.....(9)

(t,t+D) be an Where

arbitrary time interval of length D, during which i new arrivals will take place with probability

$$\frac{\left(\lambda D\right)^{i}}{i!} e^{-\lambda D}$$

,. ., . In the Erlang's integr required that F(t)=0 for every rid of this requirement by equation as:

$$\int_{\max\{0,u-(m-1)D\}}^{u} H(x,mD-u) dx + \sum_{k=0}^{\infty} \int_{u+kD}^{u+(k+1)D} H(x,mD-u) dx$$

By carrying the integration and realizing that max $\{0, u-(m-1) D\} = 0$, Except for m=1, then

$$e^{-\lambda u} \sum_{j=0}^{(m+1)c-1} Q_{(m-1)c-j-1} \frac{\left(\lambda u\right)^{j+c}}{j!} \sum_{i=0}^{c-1} \frac{\left(-1\right)^{i}}{i!(j+i+1)!(c-1-i)!} + \sum_{k=0}^{\infty} e^{-\lambda l(k+1)D+u} \sum_{j=0}^{(k+1)D-i} Q_{(k+m)c-j-1} \frac{\lambda^{j+c}}{j!} \sum_{k=0}^{c-1} \frac{\left(-1\right)^{i} \left\{(k+1)D+u\right\}^{c-l-i} D^{j+i+1}}{i!(j+i+1)!(c-1-i)!}$$

Now using
$$\sum_{i=0}^{c-1} \sum_{n=0}^{c-i-1} A_{i,n} = \sum_{n=0}^{c-1} \sum_{i=0}^{c-n-1} A_{i,n}$$

So the second term of (12) will be

$$e^{-\lambda u}\sum_{n=0}^{c-1}\frac{(\lambda u)^n}{n!}eta_n$$

Now

$$\sum_{i=0}^{c-n-1} \frac{\left(-1\right)^{i} \left(k+1\right)^{c-1-i-n}}{i! \left(j+i+1\right) \left(c-1-i-n\right)!} = \sum_{i=0}^{c-n-1} \frac{\left(-1\right)^{i}}{i! \left(j+i+1\right)} \sum_{l=0}^{c-n-1-i} \frac{k^{l}}{l! \left(c-n-1-i-e\right)!}$$
$$= j! \sum_{l=0}^{c-n-1} \frac{k^{l}}{l! \left(j+c-n-e\right)!}$$

Now

$$\begin{split} \beta_n &= \sum_{k=0}^{\infty} e^{-\lambda kD} \sum_{l=0}^{c-n-1} \frac{\left(\lambda kD\right)^l}{l!} \ Q_{(k+m)c-n-l-1} - \sum_{k=0}^{\infty} e^{-\lambda(k+1)D} \sum_{j=0}^{c-n-1} \frac{\left[\lambda \left(k+1\right)D\right]^j}{j!} Q_{(k+m+1)c-n-j-1} \\ &= \sum_{k=0}^{0} e^{-\lambda kD} \sum_{l=0}^{c-n-1} \frac{\left(\lambda kD\right)^l}{l!} \ Q_{(k+m)c-n-l-1} \\ &= Q_{mc-n-1} \end{split}$$

Now reconstruct equation (12) by putting the parts together:

$$g^{s} \text{ integral equation it is}$$
for every t<0. We can get ment by reformulating the
$$= e^{-\lambda u} \sum_{j=0}^{(m-1)c-j-1} Q_{(m-1)c-j-1} \frac{(\lambda u)^{j+c}}{(j+c)!} + e^{-\lambda u} \sum_{n=0}^{c-1} \frac{(\lambda u)^{n}}{n!} \beta_{n}$$

$$= e^{-\lambda u} \sum_{j=0}^{(m-1)c-j-1} Q_{(m-1)c-j-1} \frac{(\lambda u)^{j+c}}{(j+c)!} + e^{-\lambda u} \sum_{n=0}^{c-1} \frac{(\lambda u)^{n}}{n!} \beta_{n}$$

$$= e^{-\lambda u} \sum_{j=0}^{m-1} Q_{mc-i-1} \frac{(\lambda u)^{i}}{(j+c)!} + \sum_{n=0}^{c-1} Q_{mc-i-1} \frac{(\lambda u)^{n}}{n!} \beta_{n}$$

$$= e^{-\lambda u} \sum_{j=0}^{m-1} Q_{mc-i-1} \frac{(\lambda u)^{i}}{(j+c)!} + \sum_{n=0}^{c-1} Q_{mc-i-1} \frac{(\lambda u)^{n}}{n!} \beta_{n}$$

$$= e^{-\lambda u} \sum_{j=0}^{m-1} Q_{mc-i-1} \frac{(\lambda u)^{i}}{(j+c)!} + \sum_{n=0}^{c-1} Q_{mc-i-1} \frac{(\lambda u)^{n}}{n!} \beta_{n}$$

Now substituting y = mD-x, with m = y/D + 1, So,

This implies that the integral equation (10)

reduce to
$$F(y) = e^{-\lambda u} \sum_{i=0}^{mc-1} Q_{mc-i-1} \frac{(\lambda u)^i}{i!}$$

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